# Representing semigroups and groups by endomorphisms of Fraïssé limits 

Part II. Groups: overt \& covert

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## Sherlock Holmes: A Game of Shadows (2011)



It's so overt, it's covert - a more brutal version


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$\mathscr{H}=\mathscr{R} \cap \mathscr{L}$
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maximal subgroups of a semigroup $=\mathscr{H}$-classes containing idempotents

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Hence, $a$ is regular $\Longleftrightarrow a \mathscr{D} e$ for and idempotent $e$.

## A regular $\mathscr{D}$-class



## A regular eggbox



## A non-regular $\mathscr{D}$-class



## A non-regular eggbox



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(6) $\mathcal{T}_{X}$ is regular.

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Lemma

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(i) If $t \in T_{H}$, then $\left.t\right|_{A f}$ is an automorphism of both $\langle A f\rangle$ and $\operatorname{im}(f)$;
(ii) the mapping $\phi:\left.\rho_{t} \mapsto t\right|_{A f}$ is an embedding of $S_{H}$ into $\operatorname{Aut}(\langle A f\rangle) \cap \operatorname{Aut}(\operatorname{im}(f))$.

## So, what the heck are the images of (idempotent) endomorphisms of Fraïssé limits?

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Examples:

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Let $\mathbf{C}$ be a neat Fraïssé class enjoying the strict $A P$ and the 1 PHEP. Then there exists and (idempotent) endomorphism $f$ of $F$, the Fraïssé limit of $\mathbf{C}$, such that $A \cong \operatorname{im}(f)$ if and only if $A$ is algebraically closed in $\overline{\mathbf{C}}$.


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Let $\mathbf{K}$ be a class of $L$-structures. An $L$-structure $A$ is existentially (algebraically) closed (in $\mathbf{K}$ ) if for any primitive (positive) formula $\Phi(\mathbf{x})$ and any tuple a from $A$ we have already $A \models \Phi(\mathbf{a})$ whenever there is an extension $A^{\prime} \in \mathbf{K}$ of $A$ such that $A^{\prime} \models \Phi(\mathbf{a})$.

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- the random digraph,
- the random bipartite graph,
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## Proposition

A countable graph $(V, E)$ is a.c. if and only if there exists $E^{\prime} \subseteq E$ such that $\left(V, E^{\prime}\right) \cong R$ (that is, it is e.c.).

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## Proposition

A countable graph $(V, E)$ is a.c. if and only if there exists $E^{\prime} \subseteq E$ such that $\left(V, E^{\prime}\right) \cong R$ (that is, it is e.c.). Consequently, for any a.c. graph $\Gamma$ there is a bijective homomorphism $R \rightarrow \Gamma$.

## Frucht's Theorem (1939)

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Name of the game: Strengthen this for countable a.c. graphs.

## The team



Point Guard:
Martyn Quick


Forward:
"Baby" James Mitchell


Shooting Guard: Robert "Bob" Gray


Power Forward: Dr. D

Happy 30th birthday, Jay !!! (July 28)


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Theorem
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- $\Delta$ any graph, $\Lambda$ infinite locally finite graph $\Rightarrow(\Delta \uplus \Lambda)^{\dagger}$ is a.c.
- The central idea - consider I.f. graphs $L_{S}$ for $S \subseteq \mathbb{N} \backslash\{0,1\}$ :



## Automorphism groups of countable a.c. graphs

## Proof (cont'd).

- Properties of $L_{S}(S, T \subseteq \mathbb{N} \backslash\{0,1\})$ :
- Each $L_{S}$ is rigid $\left(\operatorname{Aut}\left(L_{S}\right)=1\right)$.
- $L_{S} \cong L_{T} \Longleftrightarrow S=T$.


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- If $L_{S}$ is isomorphic to no connected component of $\Gamma$ (and this excludes only countably many choices of $S$ ), then

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\operatorname{Aut}\left(\Gamma \uplus L_{S}\right)^{\dagger}=\operatorname{Aut}\left(\Gamma \uplus L_{S}\right) \cong \operatorname{Aut}(\Gamma) \times \operatorname{Aut}\left(L_{S}\right) \cong \operatorname{Aut}(\Gamma)
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- $S_{1} \neq S_{2}$ yield non-isomorphic a.c. graphs.


## Images of idempotent endomorphisms

Theorem (Bonato, Delić, 2000; ID, 2012)
Let $\Gamma$ be a countable graph. There exists an idempotent $f \in \operatorname{End}(R)$ such that $\operatorname{im}(f) \cong \Gamma$ if and only if $\Gamma$ is a.c.

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Theorem
If $\Gamma$ is a countable a.c. graph, then there exists an (induced) subgraph $\Gamma^{\prime} \cong \Gamma$ of $R$ such that there are $2^{\aleph_{0}}$ idempotent endomorphisms $f$ of $R$ such that $\operatorname{im}(f)=\Gamma^{\prime}$.

## Images of idempotent endomorphisms

## Proof.



At each stage of extending a homomorphism $\phi: \Gamma \rightarrow R_{\Gamma}$ to an endomorphism $\hat{\phi}$ of $R=R_{\Gamma}$, instead of mapping $v_{S} \mapsto v_{S \phi}$, if $\operatorname{im}(\phi)$ is a.c. one can find a common neighbour $w$ for $S \phi$ within $\operatorname{im}(\phi)$.

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In fact, at each stage there are infinitely many choices for $w$, which results in $\aleph_{0}^{\aleph_{0}}=2^{\aleph_{0}}$ extensions.

## The number of regular $\mathscr{D}$-classes with a given group

 $\mathscr{H}$-classTheorem
(i) Let $\Gamma$ be a countable graph. Then there exist $2^{\aleph_{0}}$ distinct regular $\mathscr{D}$-classes of $\operatorname{End}(R)$ whose group $\mathscr{H}$-classes are $\cong \operatorname{Aut}(\Gamma)$.

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Corollary
$\operatorname{End}(R)$ has $2^{\aleph_{0}}$ regular $\mathscr{D}$-classes. (You know, the ones with eggs...)

## The size of a regular eggbox

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Every regular $\mathscr{D}$-class of $\operatorname{End}(R)$ contains $2^{\aleph_{0}}$ many $\mathscr{R}$ - and $\mathscr{L}$-classes.

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All such $f$ are idempotents, and $f \mathscr{D} e$, moreover, $f \mathscr{L}$ e.
However, all these idempotents are not $\mathscr{R}$-related.

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Construct $R$ around $\Gamma^{\sharp}$, so that $R=R_{\Gamma^{\sharp}}$.
$\Gamma$ a.c. $\Longrightarrow \Gamma^{\sharp}$ a.c. Hence, the identity map on $\Gamma^{\sharp}$ can be extended to an endomorphism $g: R \rightarrow \Gamma^{\sharp}$.

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For each binary sequence $\mathbf{b}=\left(b_{i}\right)_{i \in \mathbb{N}}$ define a map $\psi_{\mathbf{b}}$ on $\Gamma^{\sharp}$ by

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v_{i, r} \psi_{\mathbf{b}}=v_{i, b_{i}}
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for all $i \in \mathbb{N}$ and $r \in\{0,1\}$.

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$g \psi_{\mathbf{b}} \in \operatorname{End}(R)$ are idempotents, $\operatorname{im}\left(g \psi_{\mathbf{b}}\right) \cong \Gamma \Rightarrow$ all these idempotents are $\mathscr{D}$-related to $e$.

Different images $\Rightarrow$ they are not $\mathscr{L}$-related.

## Non-regular eggboxes

Theorem
Let $\Gamma \not \equiv R$ be a countable a.c. graph. Then there exists a non-regular endomorphism of $R$ such that $\operatorname{im}(f) \cong \Gamma$ and $D_{f}$ contains $2^{\aleph_{0}}$ many $\mathscr{R}$ - and $\mathscr{L}$-classes.

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The proof is a variation of the idea of $\Gamma^{\sharp}$ and binary sequences.

Theorem
There are $2^{\aleph_{0}}$ non-regular $\mathscr{D}$-classes in $\operatorname{End}(R)$.
Open Problem
Are there any non-regular eggboxes of some other size?

## Schützenberger groups in $\operatorname{End}(R)$

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## Proposition

Let $f$ be an injective endomorphism of $R=(V, E)$ as described above, with $V f=V_{0}$. Then

$$
S_{H_{f}} \cong \operatorname{Aut}\left(\left\langle V_{0}\right\rangle\right) \cap \operatorname{Aut}(\operatorname{im}(f))
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## Schützenberger groups in $\operatorname{End}(R)$

So, to show a universality result for Schützenberger groups in $\operatorname{End}(R)$, one needs to extend the Frucht-de Groot-Sabidussi Theorem to countable a.c. graphs with 2-coloured edges (blue and red, say) where the 'red graph' is $\cong R$.

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This is what we did via an involved construction that again uses the rigid graphs $L_{S}$ (for a particular countable family of sets $S$ ).

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## Theorem

Let $\Gamma$ be any countable graph. There are $2^{\aleph_{0}}$ non-regular $\mathscr{D}$-classes of $\operatorname{End}(R)$ such that the Schützenberger groups of the $\mathscr{H}$-classes within them are $\cong \operatorname{Aut}(\Gamma)$.

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See arXiv:1408.4107 for details.

## A few words on posets

A poset $(P, \leq)$ is a.c. if for any finite $A, B \subseteq P$ such that $A \leq B$ there exists $x \in P$ such that

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It follows all countable/finite groups arise as automorphism groups of countable/finite a.c. posets.

## A few words on posets

However, for strict posets $(P,<)$ the notion of being a.c. changes: here we require that for all finite $A<B$ we have $x \in P$ such that

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## Open Problem

What are the automorphism groups of countable a.c. strict posets?
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Related work: G. Behrendt (PEMS, 1992)

## THANK YOU!

Questions and comments to: dockie@dmi.uns.ac.rs

Further information may be found at: http://people.dmi.uns.ac.rs/~dockie

