Representing semigroups and groups by endomorphisms of Fraïssé limits

Part II. Groups: overt & covert

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Sherlock Holmes: A Game of Shadows (2011)



It's so overt, it's covert - a more brutal version



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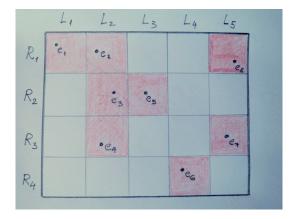
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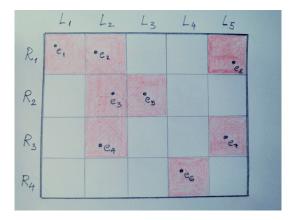
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The eggbox picture of a \mathcal{D} -class

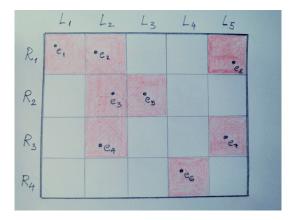


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maximal subgroups of a semigroup = \mathscr{H} -classes containing idempotents

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$$a = axa$$

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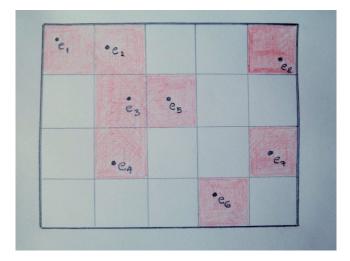
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Hence, a is regular $\iff a \mathscr{D} e$ for and idempotent e.

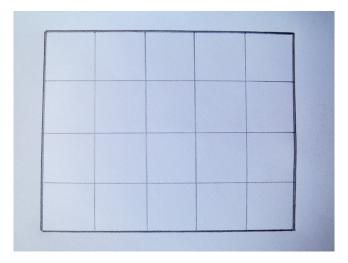
A regular \mathscr{D} -class



A regular eggbox



A non-regular \mathscr{D} -class



A non-regular eggbox



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Hence, $S_H = \{\rho_t : t \in T_H\}$ is a permutation group on H. This is the (right) Schützenberger group of H.

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Lemma

 $f \mathscr{D} g \implies \langle Af \rangle \cong \langle Ag \rangle.$

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- Let $f \in End(A)$ and $H = H_f$.
 - (i) If t ∈ T_H, then t|_{Af} is an automorphism of both ⟨Af⟩ and im(f);
- (ii) the mapping $\phi : \rho_t \mapsto t|_{Af}$ is an embedding of S_H into $Aut(\langle Af \rangle) \cap Aut(im(f))$.

So, what the heck are the images of (idempotent) endomorphisms of Fraïssé limits?

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Theorem (ID, 2012)

Let **C** be a neat Fraïssé class enjoying the strict AP and the 1PHEP. Then there exists and (idempotent) endomorphism f of F, the Fraïssé limit of **C**, such that $A \cong im(f)$ if and only if A is algebraically closed in \overline{C} .

An *L*-formula $\Phi(\mathbf{x})$ is primitive if it is of the form

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Let **K** be a class of *L*-structures. An *L*-structure *A* is existentially (algebraically) closed (in **K**) if for any primitive (positive) formula $\Phi(\mathbf{x})$ and any tuple **a** from *A* we have already $A \models \Phi(\mathbf{a})$ whenever there is an extension $A' \in \mathbf{K}$ of *A* such that $A' \models \Phi(\mathbf{a})$.

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Proposition

A countable graph (V, E) is a.c. if and only if there exists $E' \subseteq E$ such that $(V, E') \cong R$ (that is, it is e.c.). Consequently, for any a.c. graph Γ there is a bijective homomorphism $R \to \Gamma$.

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Name of the game: Strengthen this for countable a.c. graphs.

The team



Point Guard: Martyn Quick



Forward: "Baby" James Mitchell



Center: Jillian "Jay" McPhee



Shooting Guard: Robert "Bob" Gray



Power Forward: Dr. D

Happy 30th birthday, Jay !!! (July 28)



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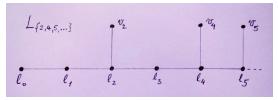
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- The central idea consider l.f. graphs L_S for $S \subseteq \mathbb{N} \setminus \{0, 1\}$:



Proof (cont'd).

- Properties of L_S ($S, T \subseteq \mathbb{N} \setminus \{0, 1\}$):
 - Each L_S is rigid $(Aut(L_S) = 1)$.

•
$$L_S \cong L_T \iff S = T$$

Proof (cont'd).

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 - Each L_S is rigid (Aut(L_S) = 1).
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- If L_S is isomorphic to no connected component of Γ (and this excludes only countably many choices of S), then

 $\operatorname{Aut}(\Gamma \uplus L_{\mathcal{S}})^{\dagger} = \operatorname{Aut}(\Gamma \uplus L_{\mathcal{S}}) \cong \operatorname{Aut}(\Gamma) \times \operatorname{Aut}(L_{\mathcal{S}}) \cong \operatorname{Aut}(\Gamma).$

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•
$$S_1 \neq S_2$$
 yield non-isomorphic a.c. graphs.

Theorem (Bonato, Delić, 2000; ID, 2012)

Let Γ be a countable graph. There exists an idempotent $f \in \text{End}(R)$ such that $\text{im}(f) \cong \Gamma$ if and only if Γ is a.c.

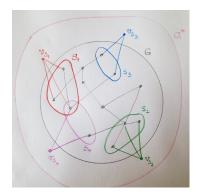
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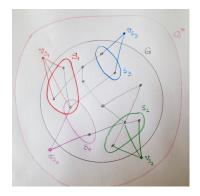
If Γ is a countable a.c. graph, then there exists an (induced) subgraph $\Gamma' \cong \Gamma$ of R such that there are 2^{\aleph_0} idempotent endomorphisms f of R such that $\operatorname{im}(f) = \Gamma'$.

Proof.



At each stage of extending a homomorphism $\phi: \Gamma \to R_{\Gamma}$ to an endomorphism $\hat{\phi}$ of $R = R_{\Gamma}$, instead of mapping $v_S \mapsto v_{S\phi}$, if $im(\phi)$ is a.c. one can find a common neighbour w for $S\phi$ within $im(\phi)$.

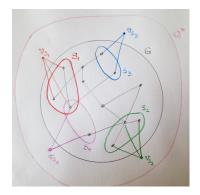
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In this way, we achieve $\operatorname{im}(\hat{\phi}) = \operatorname{im}(\phi).$

Proof.



At each stage of extending a homomorphism $\phi: \Gamma \to R_{\Gamma}$ to an endomorphism $\hat{\phi}$ of $R = R_{\Gamma}$, instead of mapping $v_S \mapsto v_{S\phi}$, if $im(\phi)$ is a.c. one can find a common neighbour w for $S\phi$ within $im(\phi)$.

In this way, we achieve $\operatorname{im}(\hat{\phi}) = \operatorname{im}(\phi).$

In fact, at each stage there are infinitely many choices for w, which results in $\aleph_0^{\aleph_0} = 2^{\aleph_0}$ extensions.

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Corollary

End(R) has 2^{\aleph_0} regular \mathcal{D} -classes. (You know, the ones with eggs...)

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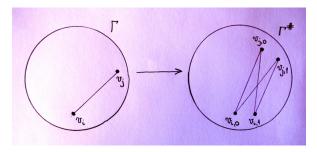
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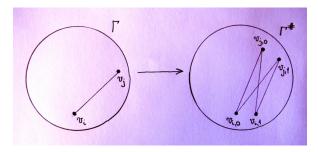
However, all these idempotents are not \mathscr{R} -related.

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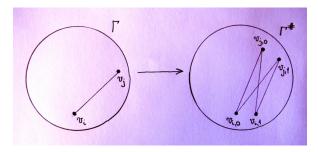


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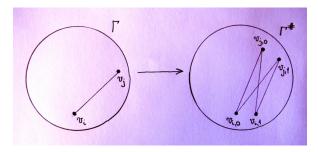
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Construct *R* around Γ^{\sharp} , so that $R = R_{\Gamma^{\sharp}}$.

 Γ a.c. $\Longrightarrow \Gamma^{\sharp}$ a.c. Hence, the identity map on Γ^{\sharp} can be extended to an endomorphism $g : R \to \Gamma^{\sharp}$.

For each binary sequence $\mathbf{b} = (b_i)_{i \in \mathbb{N}}$ define a map $\psi_{\mathbf{b}}$ on Γ^{\sharp} by

$$v_{i,r}\psi_{\mathbf{b}} = v_{i,b_i}$$

for all $i \in \mathbb{N}$ and $r \in \{0, 1\}$.

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Different images \Rightarrow they are not \mathscr{L} -related.

Theorem

Let $\Gamma \ncong R$ be a countable a.c. graph. Then there exists a non-regular endomorphism of R such that $im(f) \cong \Gamma$ and D_f contains 2^{\aleph_0} many \mathscr{R} - and \mathscr{L} -classes.

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Theorem There are 2^{\aleph_0} non-regular \mathcal{D} -classes in End(R).

Open Problem

Are there any non-regular eggboxes of some other size?

Schützenberger groups in End(R)

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Proposition

Let f be an injective endomorphism of R = (V, E) as described above, with $Vf = V_0$. Then

$$S_{H_f} \cong \operatorname{Aut}(\langle V_0 \rangle) \cap \operatorname{Aut}(\operatorname{im}(f))$$

So, to show a universality result for Schützenberger groups in End(R), one needs to extend the Frucht-de Groot-Sabidussi Theorem to countable a.c. graphs with 2-coloured edges (blue and red, say) where the 'red graph' is $\cong R$.

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Let Γ be any countable graph. There are 2^{\aleph_0} non-regular \mathcal{D} -classes of $\operatorname{End}(R)$ such that the Schützenberger groups of the \mathscr{H} -classes within them are $\cong \operatorname{Aut}(\Gamma)$.

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See arXiv:1408.4107 for details.

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It follows all countable/finite groups arise as automorphism groups of countable/finite a.c. posets.

However, for strict posets (P, <) the notion of being a.c. changes: here we require that for all finite A < B we have $x \in P$ such that

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What are the automorphism groups of countable a.c. strict posets? (I.e. what are the maximal subgroups of $End(\mathbb{P}, <)$?)

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Related work: G. Behrendt (PEMS, 1992)

THANK YOU!

Questions and comments to: dockie@dmi.uns.ac.rs

Further information may be found at: http://people.dmi.uns.ac.rs/~dockie