

One-relator groups, monoids, inverse monoids: An update on the word problem

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The word problem (in groups, monoids,...)

Assume we have given a (finitely generated) group $G = \langle X \rangle$ (e.g. by a presentation, etc.). So, elements of G are represented by **words** over $\bar{X} = X \cup X^{-1}$.

The **word problem** for G is the following decision (algorithmic) problem:

INPUT: A word $w \in \bar{X}^*$.

QUESTION: Does w represent the identity element 1 in G ?

Similarly, one can ask about the word problem for **monoids / inverse monoids / ...**, with the difference being that the input requires **two** words u, v (over X^* or \bar{X}^* , respectively), and then we want to decide if $u = v$ holds in the corresponding monoid.

One-relator groups

A **one-relator group** is a group defined by the presentation of the form

$$G = \text{Gp}\langle X \mid w = 1 \rangle = \text{FG}(X) / \langle\langle w \rangle\rangle$$

for a word $w \in \bar{X}$.

W. Magnus (1932): Every one-relator group has a decidable word problem.

Examples:

- ▶ $\mathbb{Z} \times \mathbb{Z} = \text{Gp}\langle x, y \mid [x, y] = 1 \rangle$
- ▶ Baumslag-Solitar groups
 $B(m, n) = \text{Gp}\langle a, b \mid b^{-1}a^m b a^{-n} = 1 \rangle$
- ▶ (orientable) surface groups
 $\text{Gp}\langle a_1, \dots, a_g, b_1, \dots, b_g \mid [a_1, b_1] \dots [a_g, b_g] = 1 \rangle$



So, what about one-relator monoids?

Open Problem (as of 2021 (!))

Is the word problem decidable for all one-relator monoids $\text{Mon}\langle X \mid u = v \rangle$?

Theorem (Adyan, 1966)

The word problem for $\text{Mon}\langle X \mid u = v \rangle$ is decidable if either:

- ▶ *one of u, v is empty (e.g. $u = 1$ – **special monoids**), or*
- ▶ *both u, v are non-empty, and have different initial letters and different terminal letters.*

Lallement (1977) and **L. Zhang** (1992) provided alternative proofs for the result about special monoids. (The proof of Zhang is particularly compact and elegant.)

The connection to the inverse realm

Adyan & Oganessyan (1987): The word problem for one-relator monoids can be reduced to the special case of

$$\text{Mon}\langle X \mid asb = atc \rangle$$

where $a, b, c \in X$, $b \neq c$ and $s, t \in X^*$ (and their duals).

So, where do (one-relator) **inverse** monoids come into play?

Theorem (Ivanov, Margolis & Meakin, 2001)

*If the word problem is decidable for all **special inverse monoids** $\text{Inv}\langle X \mid w = 1 \rangle$ — where w is a reduced word over \bar{X} — then the word problem is decidable for every one-relator monoid.*

This holds basically because $M = \text{Mon}\langle X \mid asb = atc \rangle$ embeds into $I = \text{Inv}\langle X \mid asbc^{-1}t^{-1}a^{-1} = 1 \rangle$.

The plot thickens

	$\text{Gp}\langle X \mid w = 1 \rangle$	$\text{Mon}\langle X \mid w = 1 \rangle$	$\text{Inv}\langle X \mid w = 1 \rangle$
decidable WP	✓ (Magnus, 1932)	✓ (Adyan, 1966)	? ✗ (Gray, 2019)

Conjecture (Margolis, Meakin, Stephen, 1987)

Every inverse monoid of the form $\text{Inv}\langle X \mid w = 1 \rangle$ has decidable word problem.

Theorem (RD Gray, *Invent. Math.*, 2020)

There exists a one-relator inverse monoid $\text{Inv}\langle X \mid w = 1 \rangle$ with undecidable word problem.

This result follows from the existence of a particular one-relator group G and its finitely generated submonoid $N \leq G$ such that the **membership problem** for N in G is undecidable.

The (prefix) membership problem (1)

Let G be a finitely generated group, and let $M \leq G$ be a finitely generated submonoid. The **membership problem for M in G** is the following decision problem:

INPUT: A word $w \in \bar{X}^*$.

QUESTION: Does the element of G represented by w belong to M ?

Now let $G = \text{Gp}\langle X \mid w = 1 \rangle$. The **prefix monoid** of G is the submonoid of G generated by all the elements represented by the prefixes of w :

$$P_w = \text{Mon}\langle \text{pref}(w) \rangle \leq G.$$

The **prefix membership problem** for the (one-relator) group G asks if the membership problem for P_w in G is decidable.

(**Caution:** depends on the presentation!)

The (prefix) membership problem (2)

Example

Let $G = \text{Gp}\langle a, b \mid aba^{-1}b^{-1} = 1 \rangle$.

$P_w = \text{Mon}\langle a, ab, aba^{-1} \rangle = \text{Mon}\langle a, b \rangle$ (because $aba^{-1} = b$ in G).

So P_w consists of elements of G represented by all positive words.

The prefix membership problem is decidable in this case, as it suffices to bring a given word into a normal form $a^i b^j$ and then check whether $i, j \geq 0$.

For example:

$$a^5 b^{-7} a^{-10} b^8 a^9 = a^4 b \in P_w, \quad a^3 b^5 a^{-5} b^2 = a^{-2} b^7 \notin P_w.$$

Problem

Does every one-relator group $G = \text{Gp}\langle X \mid w = 1 \rangle$ has decidable prefix membership problem? If the answer is “no”, characterise words w which do have this property.

The (prefix) membership problem (3)

Proved true in several special cases, when $w \dots$

- ▶ \dots represents an idempotent of the free inverse monoid (Birget, Margolis, Meakin, 1993/4);
- ▶ \dots is “strictly positive” (Ivanov, Margolis, Meakin, 2001);
- ▶ \dots defines certain Adyan-type or Baumslag-Solitar-type groups (Margolis, Meakin, Šunić, 2005);
- ▶ \dots satisfies certain small-cancellation-type conditions (Juhász, 2012, 2014).

Tying everything together

Theorem (Ivanov, Margolis & Meakin, 2001)

If $M = \text{Inv}\langle X \mid w = 1 \rangle$ is E -unitary, then

word problem for M = prefix membership problem for $G = \text{Gp}\langle X \mid w = 1 \rangle$.

E-unitary inverse semigroups = the “good guys” of inverse semigroup theory:

- ▶ For any $e \in E(S)$ and $x \in S$,
 $e \leq x$ (in the natural inverse semigroup order) $\Rightarrow x \in E(S)$.
- ▶ The minimum group congruence σ on S is **idempotent-pure**, meaning that $E(S)$ constitutes a single σ -class (the identity element of the group S/σ).
- ▶ ...

Theorem (Ivanov, Margolis & Meakin, 2001)

If w is **cyclically reduced**, then $M = \text{Inv}\langle X \mid w = 1 \rangle$ is E -unitary.

The paper

In our recently published paper,

I.Dolinka, R.D.Gray, New results on the prefix membership problem for one-relator groups, *Trans. Amer. Math. Soc.* **374** (2021), 4309–4358.

our strategy was first to prove several general sufficient conditions for a finitely generated submonoid in

- ▶ free amalgamated products of f.g. groups,
- ▶ HNN-extensions of a f.g. group,

and then to apply these results to deduce decidability of the prefix membership problem in certain classes of one-relator groups.

Here is a sample of such a type of result.

Theorem A

Let $G = H *_L K$, where H, K, L are finitely generated groups such that both H, K have decidable word problems, and the membership problem for L in both H and K is decidable.

Let M be a submonoid of G such that:

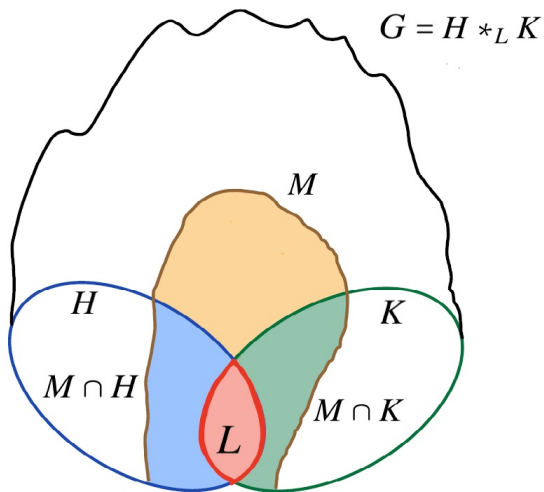
- (i) $L \subseteq M$;
- (ii) both $M \cap H$ and $M \cap K$ are finitely generated, and

$$M = \text{Mon}\langle (M \cap H) \cup (M \cap K) \rangle;$$

- (iii) the membership problem for $M \cap H$ in H is decidable;
- (iv) the membership problem for $M \cap K$ in K is decidable.

Then the membership problem for M in G is decidable.

Picture for Theorem A



Conservative factorisations (1)

Let $G = \text{Gp}\langle X \mid w = 1 \rangle$ and consider a factorisation

$$w \equiv w_1 \dots w_k.$$

Define $P(w_1, \dots, w_k)$ to be the submonoid of G generated by $\bigcup_{1 \leq i \leq k} \text{pref}(w_i)$.

Since every prefix of w is a product of prefixes of w_i 's, clearly $P_w \subseteq P(w_1, \dots, w_k)$. When equality takes place, we say that the considered factorisation is **conservative**.

The factors w_i in a conservative factorisation are called **pieces**.

Conservative factorisations (2)

Methods for finding conservative factorisations:

- ▶ **Adyan overlap algorithm**: based on the fact that if $\alpha\beta$ and $\beta\gamma$ are pieces, so are α, β, γ
 - ▶ **Example**: w.r.t. $G = \text{Gp}\langle a, b, c, d \mid abcdabcdcd = 1 \rangle$, the factorisation $(ab)(cd)(ab)(cd)(cd)$ is conservative
- ▶ **Benois method** [Gray & Ruškuc, 2021]: based on Benois' Theorem (stating that rational subsets in f.g. free groups have uniformly decidable membership problem)
 - ▶ **Example**:
 $G = \text{Gp}\langle a, b, c, d \mid (abcd)(acd)(ad)(abbcd)(acd) = 1 \rangle$ is a conservative factorisation

Application 1: Unique marker letter theorem

Let $G = \text{Gp}\langle X \mid w = 1 \rangle$, where $w \equiv w_1 \dots w_k$ is a conservative factorisation. Let $U = \{u_1, \dots, u_m\}$ be the set of the pieces appearing in this factorisation. Suppose that

- ▶ for all $1 \leq i \leq m$ there is a letter $a_i \in X$ appearing exactly once in u_i and not appearing in any u_j , $j \neq i$.

Then $G = \text{Gp}\langle X \mid w = 1 \rangle$ has decidable prefix membership problem.

Example

In $G = \text{Gp}\langle a, b, x, y \mid axbaybaybaxbaybaxb = 1 \rangle$, the Adyan overlap method produces a conservative factorisation

$$(axb)(ayb)(ayb)(axb)(ayb)(axb),$$

where the pieces axb and ayb have the unique marker letter property, so G has decidable prefix membership problem.

Application 1: Unique marker letter theorem

$G = \text{Gp}\langle a, b, c, d \mid (abcd)(acd)(ad)(abbc)(acd) = 1 \rangle$ is a conservative factorisation. Here, the pieces don't have the unique marker letter property. However, upon transforming

$$\begin{aligned} G &= \text{Gp}\langle a, b, c, d \mid (abcd)(acd)(ad)(abbc)(acd) = 1 \rangle \\ &= \text{Gp}\langle a, b, c, d \mid (aba^{-1})(aca^{-1})(ad)(aca^{-1})(ad)(ad) \\ &\quad (aba^{-1})(aba^{-1})(aca^{-1})(ad)(aca^{-1})(ad) = 1 \rangle \end{aligned}$$

we obtain a conservative factorisation, where the pieces aba^{-1} , aca^{-1} and ad do have the unique marker letter property.

Since the prefix monoid turns out to be unchanged in this modified presentation of G , we conclude that G (w.r.t. the initial presentation) has decidable prefix membership problem.

(This answers a question harking back to the 1987 paper of Margolis, Meakin and Stephen.)

Application 2: Disjoint alphabets theorem

Let $G = \text{Gp}\langle X \mid w = 1 \rangle$, where w is cyclically reduced and $w \equiv w_1 \dots w_k$ is a conservative factorisation. Let $U = \{u_1, \dots, u_m\}$ be the set of the pieces appearing in this factorisation. Suppose that $m \geq 2$ and

- ▶ for all $1 \leq i \neq j \leq m$, u_i and u_j have no letters in common.

Then $G = \text{Gp}\langle X \mid w = 1 \rangle$ has decidable prefix membership problem.

Example

In $G = \text{Gp}\langle a, b, c, d \mid ababcdcdababcdcdcdcdabab = 1 \rangle$, the Adyan overlap method produces a conservative factorisation

$$(abab)(cdcd)(abab)(cdcd)(cdcd)(abab),$$

where the pieces $abab$ and $cdcd$ have no letters in common. So, G has decidable prefix membership problem.

Application 3: Exponent sum zero theorem

Let $G = \text{Gp}\langle X \mid w = 1 \rangle$ where $t \in X$ and w (containing t) has t -exponent sum 0. (E.g. $w \equiv atat^2a^2t^{-3}$.)

Then (by the observation of **Moldavanskiĭ (1967)**), the following exists:

- ▶ a one-relator group $G' = \text{Gp}\langle X' \mid w' = 1 \rangle$ with $|w'| < |w|$;
- ▶ subsets $A, B \subset X'$ forming bases of free subgroups F_1, F_2 of G' ;
- ▶ an isomorphism $\phi : F_1 \rightarrow F_2$;

such that G is a HNN extension of G' w.r.t. ϕ .

Theorem

With the above notation, if G' is a free group and w is prefix t -positive, then G has decidable prefix membership problem.

Further applications

- ▶ **Cyclically pinched groups:** $\text{Gp}\langle X, Y \mid uv^{-1} = 1 \rangle$ where $u \in \overline{X}^*$ and $v \in \overline{Y}^*$

- ▶ These include the orientable surface groups

$$\text{Gp}\langle a_1, \dots, a_n, b_1, \dots, b_n \mid [a_1, b_1] \dots [a_n, b_n] = 1 \rangle$$

- ▶ and the non-orientable surface groups

$$\text{Gp}\langle a_1, \dots, a_n \mid a_1^2 \dots a_n^2 = 1 \rangle$$

- ▶ **Conjugacy pinched groups:** $\text{Gp}\langle X, t \mid t^{-1}utv^{-1} = 1 \rangle$ where $u, v \in \overline{X}^*$ are non-empty and reduced (these include the Baumslag-Solitar groups)
- ▶ Some **Adyan-type groups:** $\text{Gp}\langle X \mid uv^{-1} = 1 \rangle$, $u, v \in X^*$ are positive words such that the first letters of u, v are different and also the last letters of u, v are different (some new cases are covered)

A negative result and a problem

Theorem

*There exists a finite alphabet X and a **reduced** word $w \in \overline{X}^*$ such that $G = \text{Gp}\langle X \mid w = 1 \rangle$ has undecidable prefix membership problem.*

Open Problem

Characterise the words w such that $G = \text{Gp}\langle X \mid w = 1 \rangle$ has decidable prefix membership problem. Are all **cyclically reduced** words among them?

Thank you!

Questions and comments to:

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Further information may be found at:

<http://people.dmf.uns.ac.rs/~dockie>