



# **A Nonfinitely Based Finite Semiring**

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# The finite basis problem

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$\mathbf{A}$  – a finite algebra

$\text{Eq}(\mathbf{A})$  – the set of all identities true in  $\mathbf{A}$

Is  $\text{Eq}(\mathbf{A})$  finitely axiomatizable  
(finitely based)?

McKenzie (1996): in general, undecidable



# Finitely based finite algebras

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- groups: Oates & Powell (1966)
- commutative semigroups: Perkins (1968)
- lattices (& other lattice-based algebras): McKenzie (1970)
- rings: Львов, Kruse (1973)



# Some NFB finite algebras

- Мурский (1965): a 3-element groupoid
  - this is a special case of NFB graph algebras – Baker, McNulty, Werner (1987)
- Perkins (1968): a 6-element semigroup = the Brandt monoid  $B_2^1$  of order 2

$$\begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

- the Perkins' semigroup is INFB = each l.f. variety containing it is NFB (Sapir, 1987)



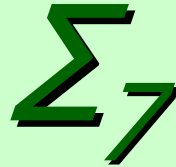
# Semirings

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**Semiring** = an algebra  $(\Sigma, +, \cdot, 0)$  such that

- $(\Sigma, +, 0)$  is a commutative monoid,
- $(\Sigma, \cdot)$  is a semigroup,
- the multiplication distributes over addition.

If  $+$  is an idempotent operation ( $x+x=x$ ),  
then we have **ai-semirings**.

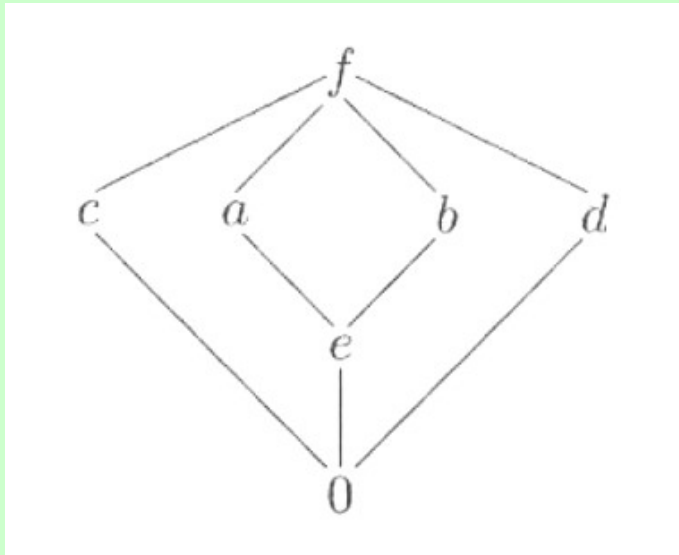


- a subsemiring of  $\mathbf{Rel}(2)$ , the semiring of binary relations on a two element set, formed by:
  - the four relations with 3 pairs,
  - the empty, the diagonal, and the full relation
- alternatively, the ai-semiring formed by 7 Boolean matrices

$$\begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}, \\ \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}, \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}.$$

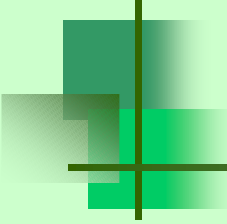
(remember that we have  $1+1=1$  in the 2-element Boolean semiring)

# $\Sigma_7$ (continued)



.	0	a	b	c	d	e	f
0	0	0	0	0	0	0	0
a	0	a	f	c	f	a	f
b	0	f	b	f	d	b	f
c	0	f	c	f	a	c	f
d	0	d	f	b	f	d	f
e	0	a	b	c	d	e	f
f	0	f	f	f	f	f	f

equations of  $B_2^1 =$  semigroup equations of  $\Sigma_7$



Is there such a thing as a NFB  
finite semiring?

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**Theorem A.**  $\Sigma_7$  is NFB.

According to *MathSciNet*, this is a first  
example of such kind.

What follows is a (hopefully) **VERY** short  
outline of the proof idea.





# IMAGIGAM words

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- a word of the form

$$yLyL^R$$

where  $L$  is a linear word not containing  $y$ , and  $L^R$  is the reverse of  $L$

- for all  $n$ ,  $B_2^1$  (and so  $\Sigma_7$ ) satisfies the imagigam equations

$$yX_1X_2 \cdots X_n yX_n \cdots X_2X_1 = yX_n \cdots X_2X_1 yX_1X_2 \cdots X_n$$



# Isoterms #1

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A word  $u$  is an isoterm for an ai-semiring identity

$$\sum_i u_i = \sum_j v_j$$

if for each semigroup substitution  $\varphi$  such that  $\varphi(u_i)$  is (for some  $i$ ) a subword of  $u$  we have that

- either not all  $\varphi$ -values of  $u_i$ 's are equal, or
- all  $\varphi$ -values of both  $u_i$ 's and  $v_j$ 's are equal



## Isoterms #2

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- for a fixed ai-semiring  $\Sigma$  and words  $u, v$  we write  $u \leq v$  if  $\Sigma$  satisfies  $u + v = v$
- a word  $w$  is minimal if  $u \leq w$  implies that  $u$  is either 0, or  $w$
- a minimal word = an isoterm for all identities of  $\Sigma$  (an isoterm of  $\Sigma$ )



## Isoterms #3

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Let  $n$  be a natural number and  $\Sigma$  an ai-semiring.

A word  $u$  in at least  $n$  letters is an  $n$ -isoterm of  $\Sigma$  if it is an isoterm for all equations of  $\Sigma$  in less than  $n$  letters.



# Why isoterms?

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**An easy proposition.** Let  $\Sigma$  be an ai-semiring. Suppose that for arbitrary large  $n$  we manage to find a word  $w_n$  which **is** an  $n$ -isoterm, but **not** an isoterm of  $\Sigma$ .

**Then  $\Sigma$  is NFB.**



# Why isoterm?

If one translates all notions to semigroups this is exactly the tool used by Perkins!

Namely, the imagigam words turn out to be suitable: Perkins proves that

$$YX_1X_2 \cdots X_n YX_n \cdots X_2X_1$$

is always a (semigroup)  $n$ -isoterm, while the imagigam equations show that it is not an isoterm of the Perkins' monoid.



# René, I've got a plan...

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Can we do the same for  $\Sigma_7$ ?

I.e., is the  $n$ th imagigam word an  $n$ -isoterm (in the ai-semiring sense) of  $\Sigma_7$ ?  
(It is obviously not an isoterm of  $\Sigma_7$ .)

How to find  $n$ -isoterms at all?

# A good lemma always saves the day!

**Lemma.** Let  $w$  be a word, with precisely  $n$  letters occurring in it, let  $\Sigma$  be an ai-semiring, and let  $k < n$  be such that

- (1) each word  $u$  in less than  $n$  letters, such that  $w$  contains a value of  $u$  (under some substitution), is minimal with respect to  $\Sigma$ ,
- (2)  $w$  satisfies a certain combinatorial (and technical – but not too much) condition called the  $k$ -joint substitution property.

Then  $w$  is a  $(k+1)$ -isoterm of  $\Sigma$ .





In  $\Sigma_n$  the imagigam words satisfy  
both conditions!

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- 1) Each word in at most  $n$  variables that has a value in the  $n$ th imagigam word is minimal in  $\Sigma_n$ .
- 2) Each imagigam word containing at least  $4k+2$  letters has the  $k$ -joint substitution property.

1) is a classical **combinatorics-on-words** issue;  
for the proof of 2) the key thing is to use a fact  
from **elementary geometry** (!)

**1) + 2) + Easy Prop. => Theorem A.**

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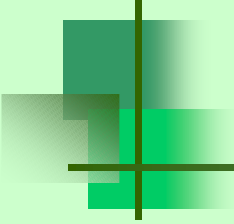
To tell the truth, we do not need the  
`full strength' of  $\text{Eq}(\Sigma_7)$ , only 7 its  
particular features so that we obtain a  
slightly more general result...

# Theorem B.

**Theorem B.** *Let  $\Sigma$  be an ai-semiring. Call  $\Sigma$  special if it satisfies the following conditions:*

- (a) *the inequalities of  $\Sigma$  are closed under deletion, i.e. for any words  $u, v$  such that  $u \preceq v$  we have  $c(u) = c(v)$ , and if  $u', v'$  are obtained respectively from  $u, v$  by deleting all occurrences of a given variable (provided  $u, v$  contain at least two variables), then  $u' \preceq v'$ ,*
- (b)  *$yx \not\preceq xy$ ,*
- (c)  *$x$  and  $xyx$  are minimal with respect to  $\Sigma$ ,*
- (d)  *$x^2y, xyx, yx^2$  are mutually  $\preceq$ -incomparable,*
- (e)  *$w \not\preceq (xy)^2$  whenever  $w \in \{xyx, yxyxy\}$  or  $w$  contains one of  $x^2, y^2$  as a subword,*
- (f)  *$xyzxy \not\preceq xyzyx, yxzyx \not\preceq xyzyx$  and  $xzyxy \not\preceq xzy^2x$ ,*
- (g)  *$w \not\preceq xyztxtz$  for  $w \in \{xytztzt, xyztxtz, xytzxzt\}$ .*

*If  $\Sigma$  is special and satisfies all the imagigam identities, then it is nonfinitely based.*



# Open questions

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- **Q1:** Are the semirings  $Rel(n)$  of binary relations on an  $n$ -element set,  $n > 1$ , finitely based or not?

- **Q2:** Is  $\Sigma_7$  INFB?

Clearly enough, A2: **Yes**  $\Rightarrow$  A1: **They're not.**

- **Q3:** If A2 is **Yes**, is the same conclusion true for each finite ai-semiring in which all Zimin words are minimal (a feature easily proved in  $\Sigma_7$  by induction)?



**Thank you!**